

# Robotics

## Robot Navigation (3)

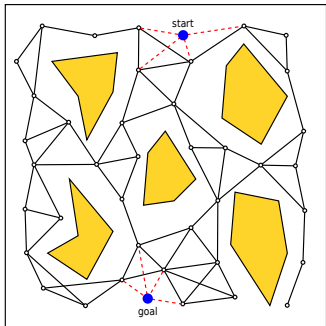
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Tuesday 3<sup>rd</sup> October, 2023

<http://robot.unipv.it/toolleeo>

# Roadmap

a **roadmap  $M$**  is a sequence of **unidimensional curves** connecting a starting point  $p_{\text{start}}$  and a goal  $p_{\text{goal}}$



## properties

- ① **accessibility**: a path exists from  $p_{\text{start}}$  to some point  $p'_s \in M$
- ② **departability**: a path exists from a point  $p'_g \in M$  to the goal point  $p_{\text{goal}}$
- ③ **connectivity**: a path exists between  $p'_s$  and  $p'_g$  made of nodes belonging to  $M$

## Probabilistic methods

when the dimension of the configuration space becomes high, some approaches are needed to **reduce the complexity of the search**: this is the objective of probabilistic methods

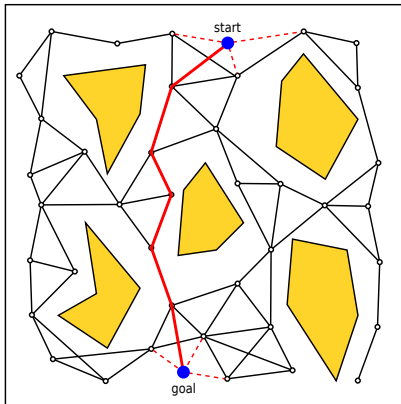
- we will consider the so-called **“sampling-based methods”**
- the father of all sampling-based methods is the **Probabilistic RoadMap (PRM)** trajectory planner

## PRM: Probabilistic RoadMap

it leverages the idea that **it is effortless** to decide whether a point belongs to the free configuration space or not

- first, a **map of the free space** is built
- the planner randomly generates points in the configuration space
- it keeps only those point that belong to the free space
- **rough generation** of point
- **accurate planning of trajectories** between collected points

# Probabilistic methods: example



# Operations of a probabilistic planner

## learning phase

- a **map** of the configuration space is built
- a **graph** is generated to connect points in the free space

## query phase

- check whether  $p_{\text{start}}$  and  $p_{\text{goal}}$  **can be connected to the map**, to nodes  $p'_s$  and  $p'_g$  respectively
- generate the **shortest path** between  $p'_s$  and  $p'_g$  (e.g. using  $A^*$ )

## Learning phase

- 1 a set of  $V$  points in the configuration space is generated
- 2 for each point  $q \in V$ , the  $k$  closest points are considered, composing the  $V_k$  set
- 3 for each point  $p \in V_k$  a path is found to connect  $p$  and  $q$ , if  $p$  and  $q$  are not yet connected

### key issues

- how the points in  $V$  are generated (randomly, using some models of obstacles, addressing the narrow passage problem, etc.)
- considering the distance between points in  $V_k$  (the computation time increases with such a distance, since more points are involved)
- how to connect pairs of points (“local planner”: straight lines, splines, etc.)

## Query phase

for both the starting point  $p_{\text{start}}$  and the goal  $p_{\text{goal}}$ :

- 1 the  $k$  closest points of the map to  $p_{\text{start}}$  and  $p_{\text{goal}}$  are considered
- 2 the possibility to connect  $p_{\text{start}}$  to  $p'_s$  and  $p_{\text{goal}}$  to  $p'_g$  is checked
- 3 the shortest path between  $p'_s$  and  $p'_g$  is searched

### key issues

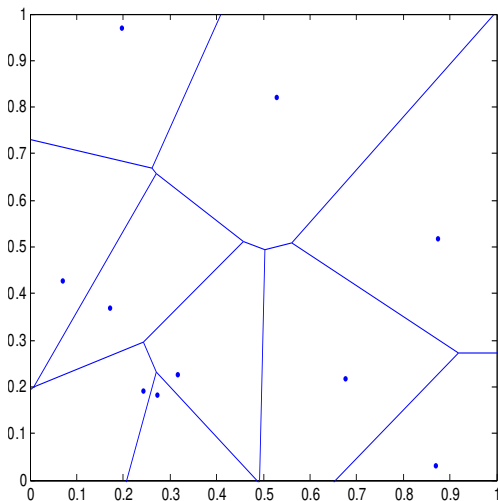
- how many nodes  $k$  shall be considered
- how to connect pairs of points (“local planner”: straight lines, splines, etc.)



# Voronoi diagrams

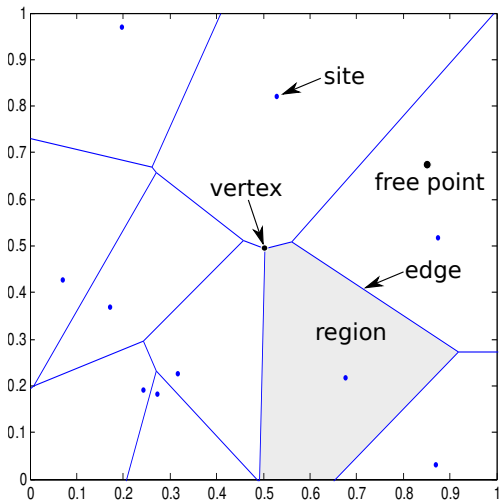
also known as:

- Voronoi tessellation
- Voronoi decomposition
- Voronoi partition
- Dirichlet tessellation

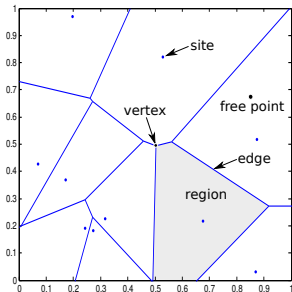


## Example and notation

- $p_i$  : sites
- $q$  : free points
- $e$  : (Voronoi) edge
- $v$  : (Voronoi) vertex
- $F_i$  : region or cell



## Voronoi region



set of sites:

$$P = \{p_i\}, 1 \leq i \leq n$$

Voronoi diagram of  $Vor(P)$ :

- division of the plane into  $n$  regions
- there is a region for each site

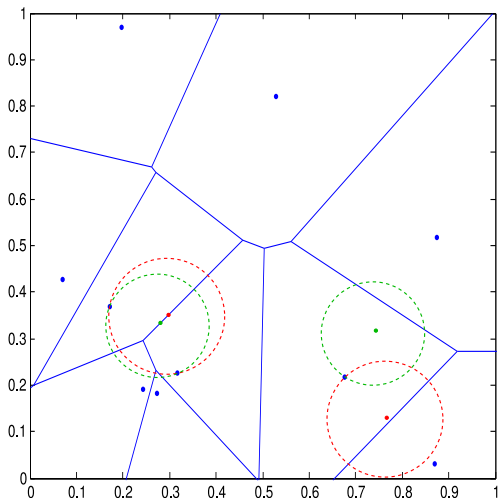
definition of the  $i$ -th region:

$$F_i = \{q \in Q_{\text{free}} : d(p_i, q) \leq d(p_j, q) \forall i \neq j\}$$

Voronoi regions are bounded by line segments

## Features of vertices

- 1 the locus of the center of a largest empty circle passing through only a pair of points  $p_i, p_j \in P$  defines an edge
- 2 the locus of the center of largest empty circles passing through only one point in  $P$  defines a region



## Construction of a Voronoi map

### sensor-based

- uses the robot onboard sensor to build the map
- the map can be built online

### polygonal spaces

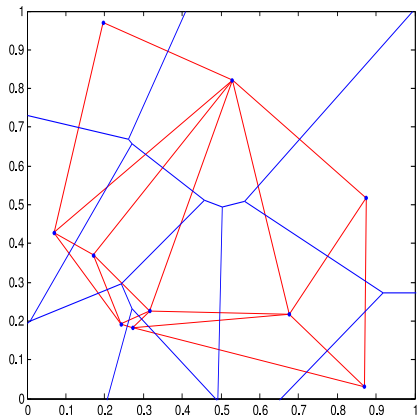
- a full representation of the environment is needed
- obstacles must be polygons

### grid map

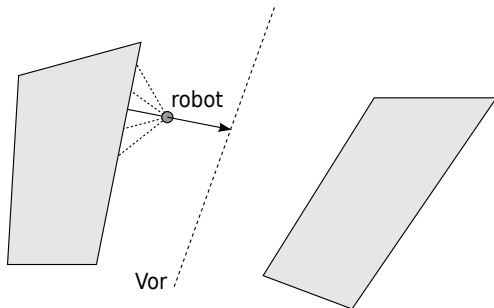
- a full representation of the environment is needed
- easy and fast to compute

## Delaunay triangulation

- Delaunay triangulation is complementary to the Voronoi tessellation
- each triangle's edge is orthogonal to a Voronoi edge
- it can be leveraged to construct the Voronoi partition
  - an algorithm is applied to build the Delaunay triangulation
  - the Voronoi graph is derived from the triangulation

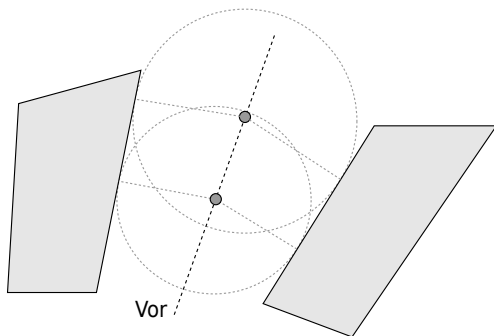


## Sensor-based construction



- the robot reaches the  $Vor$  with gradient ascent
- motion along the direction that maximise  $\nabla d(q_i, q)$  from the closest obstacle

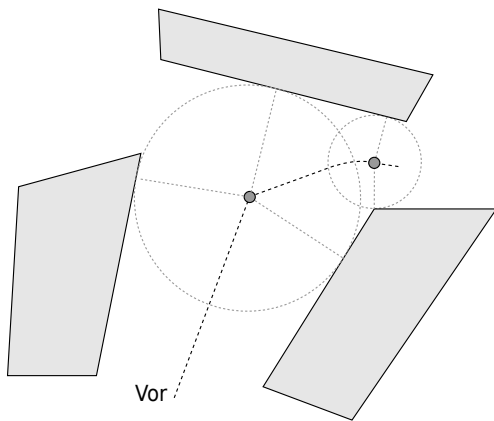
## Sensor-based construction



- the robot moves along a line that keep  $d(q_i, q) = d(q_j, q)$
- $q_i$  and  $q_j$  are the closest points belonging to the closest obstacles

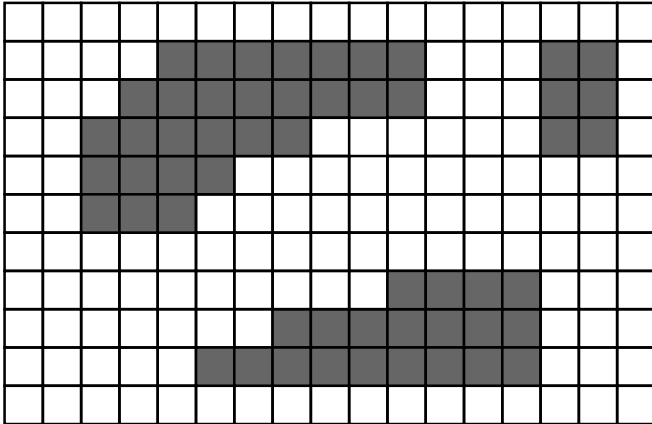


## Sensor-based construction



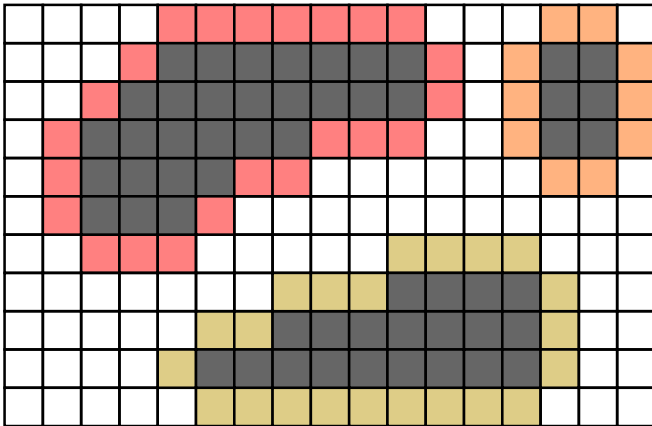
- meet points are detected from changes in the closest obstacle

## Grid-based construction



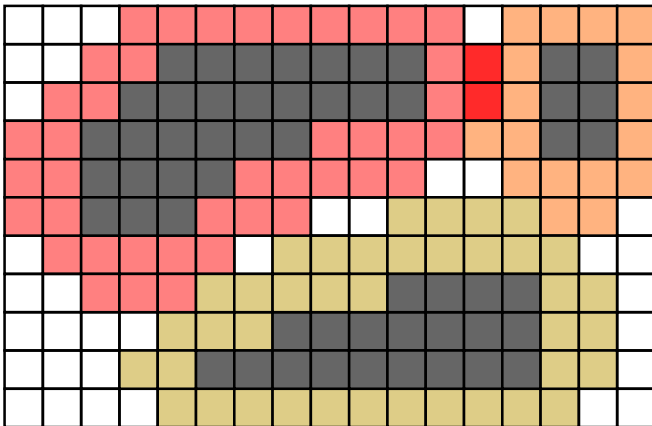
- the space is represented by means of a grid

## Grid-based construction



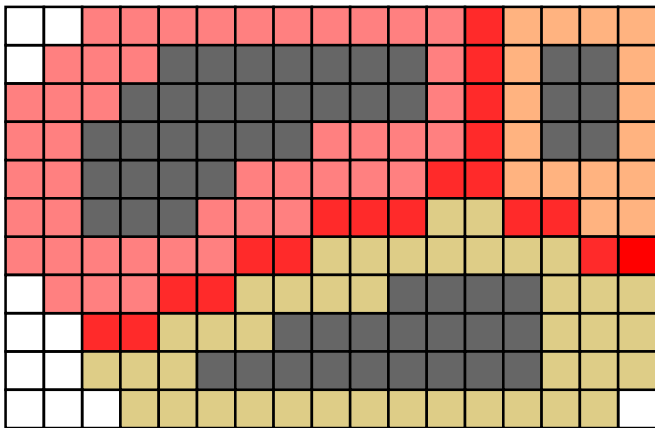
- a wave-front grow is started from each obstacle

## Grid-based construction



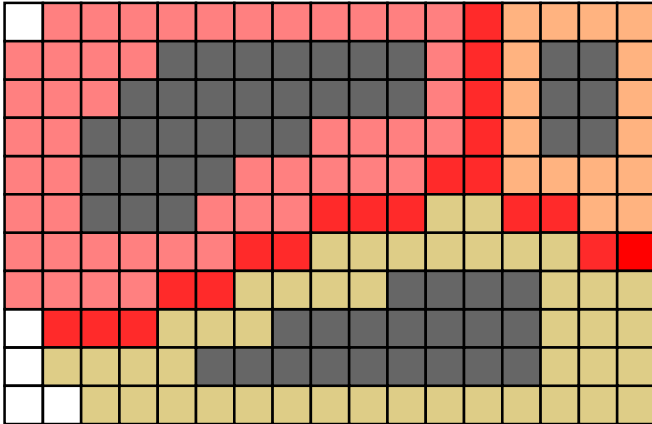
- a collision between two waves detects the presence of a *Vor* edge

## Grid-based construction



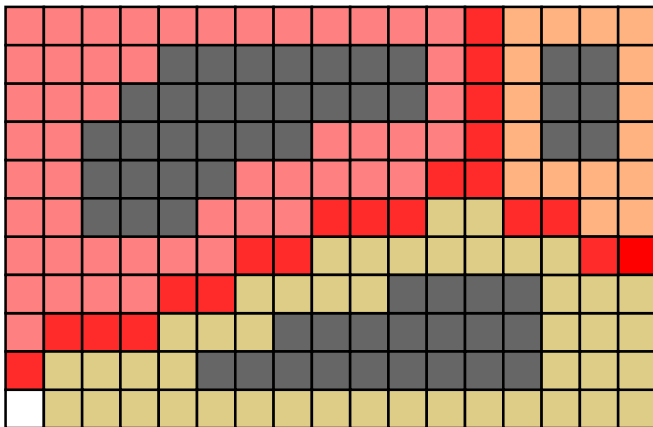
- where there is a collision “on the boundary”, the selection of the grid cell to mark is up to the algorithm

# Grid-based construction



- the algorithm continues until...

## Grid-based construction

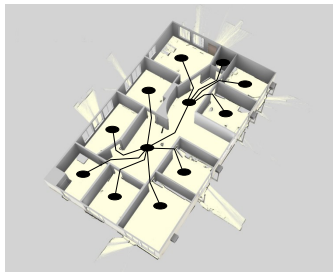


- ...the grid is fully divided into Voronoi regions

# Cell decomposition

## exact cell decomposition

the free space is represented by the union of adjacent cells (regions)



- cell boundaries are often associated to physical features of the free space
  - change in the line of sight
  - change in the closest obstacle
  - intersection with vertices
- adjacent cells share a common boundary



## Cell decomposition

### adjacency graph

it encapsulates the relationships between adjacent cells

- a node is associated to a cell or an edge (boundary)
- graph edges connect adjacent cells

### path planning

- the planner finds the cell containing the start and goal points
- the adjacency graph is used as a roadmap

# Trapezoidal decomposition

- made by two-dimensional cells having trapezoidal shape
- triangular cells are possible
  - they can be seen as degenerated trapezoid
  - one of the parallel edges is collapsed into a 0-length edge

## assumptions

- planar configuration space
- obstacles are polygonal
- the free space is bounded by a polygon

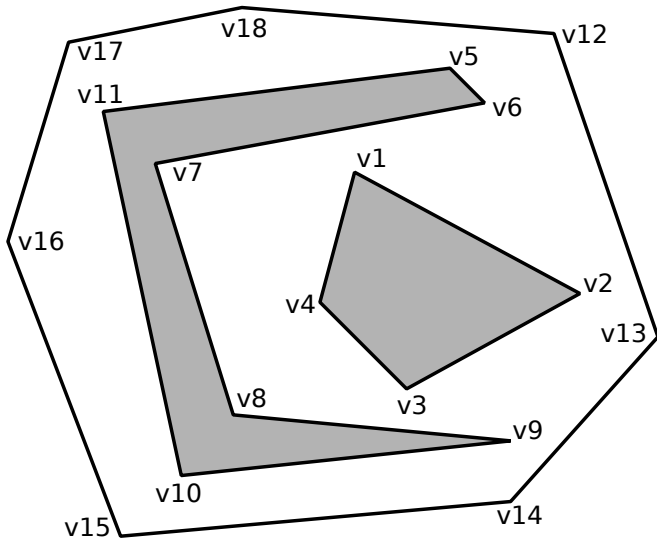
## Build the trapezoidal decomposition

- for each vertex  $v_i$ , perform the following procedure
- draw two segments originating from  $v_i$
- the extension of each segment stops when it first intersects an edge
- segments can be drawn in any direction, but they must be parallel

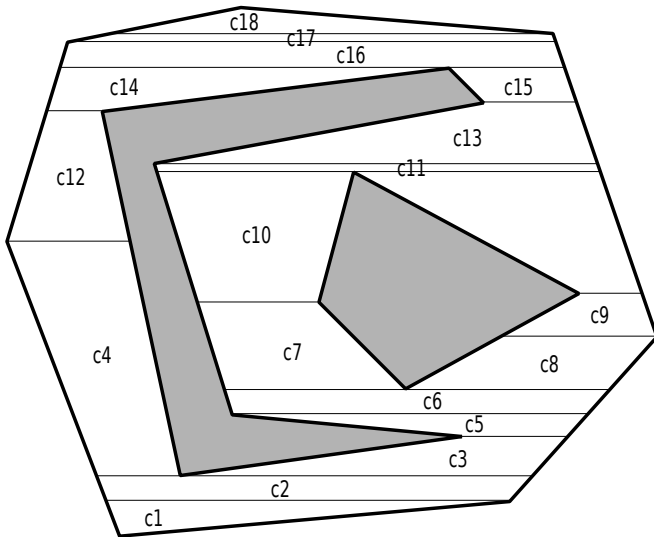
### note:

- some vertices may generate only one or even no segment

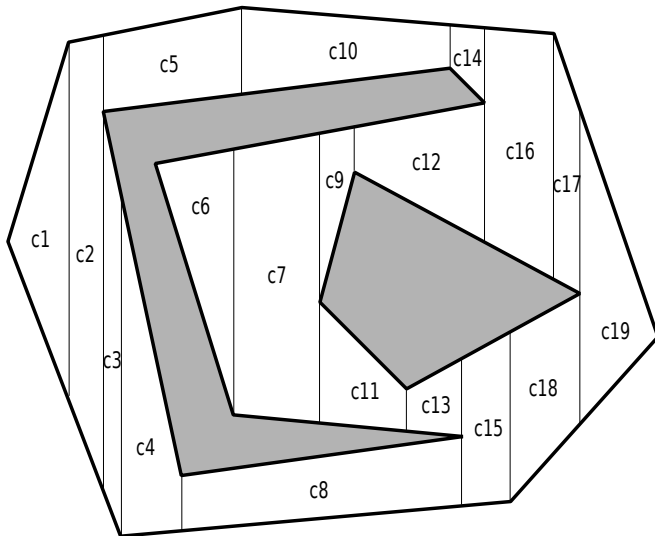
## Build the trapezoidal decomposition: example



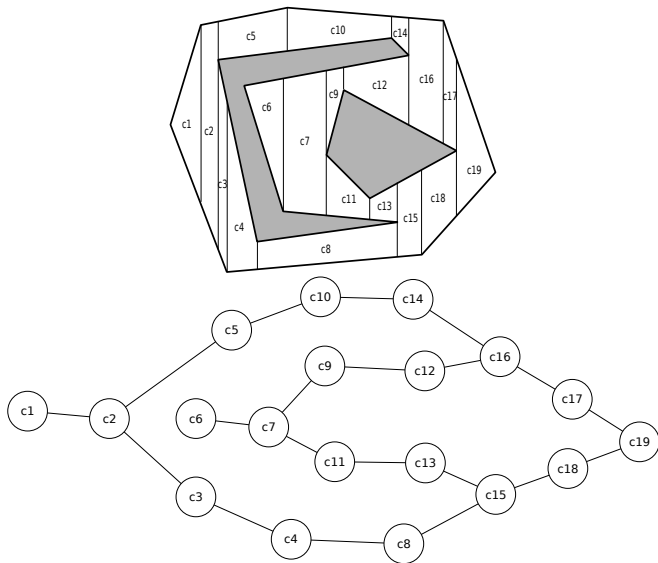
# Build the trapezoidal decomposition: example



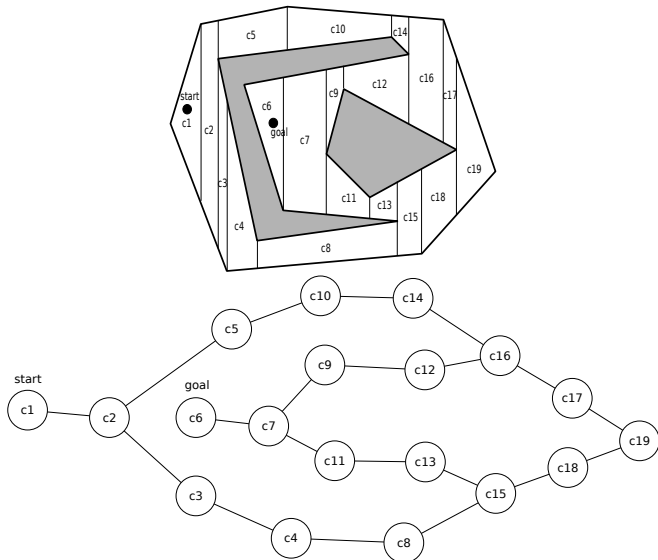
# Build the trapezoidal decomposition: example



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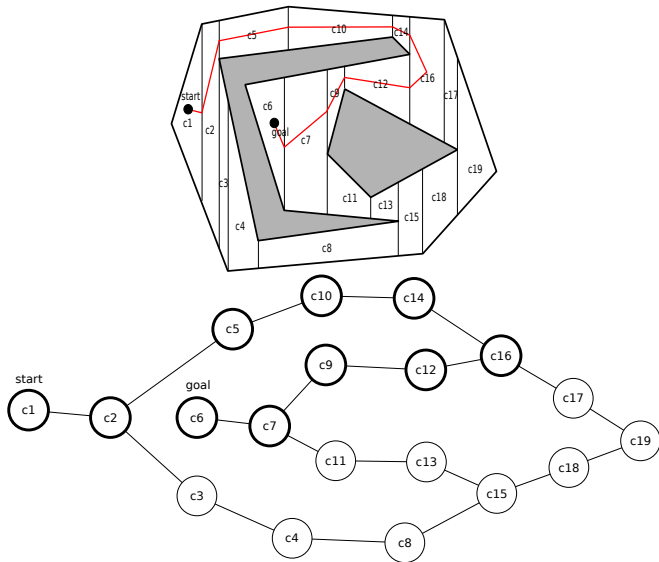


# Build the trapezoidal decomposition: example





# Build the trapezoidal decomposition: example



## Path planning

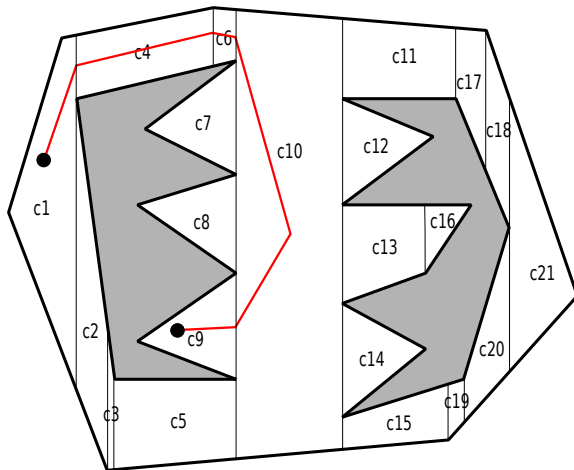
- the cells containing the start and goal points are identified
- the adjacency graph is searched for a path connecting the two cells
- the sequence of nodes in the graph must be translated into waypoints in the free space

## Path planning

trapezoids are convex  
therefore the centroid can be connected with every  
boundary segment by a straight line

- 1 beginning from the start point...
- 2 each point is connected with the mid-point of the segment associated with the next graph node
- 3 the mid-point is connected to the centroid of the trapezoid
- 4 repeat from step 2 until the cell containing the goal point is reached
- 5 connect the last mid-point with the goal point

## Properties of trapezoid



- many boundary segments are related with cell c10
- the approach to reach the centroid allows to easily reach every boundary segment

## Coverage based on cell decomposition

### coverage

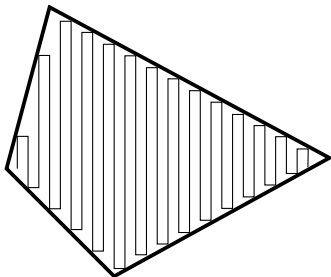
the ability to pass a sensor or effector over all points composing the free space

### applications:

- exploration
- surveillance
- manufacturing (e.g., painting, polishing, smoothing)

## Coverage based on cell decomposition

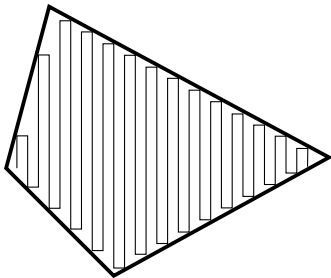
- in general, the structure of a cell is simple
- after the free space is divided into cells, simple algorithms can be used to cover each cell
- once all cells are visited, coverage is obtained



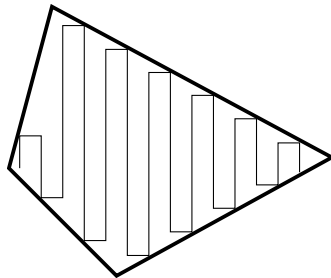
back-and-forth coverage of a cell

## Coverage based on cell decomposition

- the horizontal offset determines the trade-off between coverage accuracy and execution time
- wider offset leads to quicker but more inaccurate coverage
- shorter offset leads to more accurate coverage but takes more time



shorter offset



wider offset